

Markov chains and applications

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Week 13, University of Debrecen



Infinite vectors, matrices

$$\begin{pmatrix} a_{1,1} & a_{1,2} & a_{1,3} & \cdots \\ a_{2,1} & a_{2,2} & a_{2,3} & \cdots \\ a_{3,1} & a_{3,2} & a_{3,3} & \cdots \\ \vdots & \vdots & \vdots & \ddots \end{pmatrix}$$

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We can multiply infinite row and column vectors (indexed by \mathbb{N} or $\mathbb{N} \cup \{0\}$) in the logical way, by computing an infinite sum.

Perron-Frobenius theorem



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Perron-Frobenius theorem



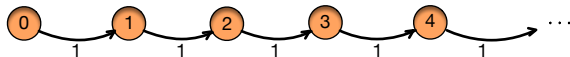
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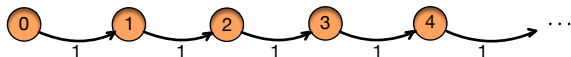


$$P = \begin{pmatrix} 0 & 1 & 0 & 0 & \dots \\ 0 & 0 & 1 & 0 & \dots \\ 0 & 0 & 0 & 1 & \dots \\ \vdots & \vdots & \vdots & \vdots & \ddots \end{pmatrix}$$

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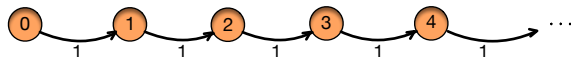


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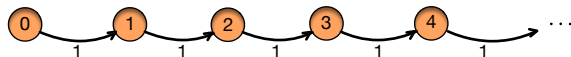
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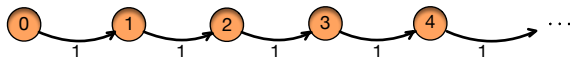
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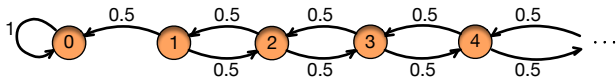
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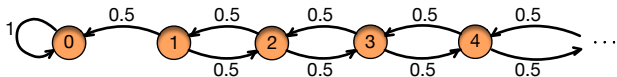
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$\lim_{k \rightarrow \infty} P^k = 0$ (the walk fades away at infinity)

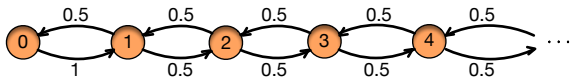
An absorbing example (infinite fair gambler's ruin):



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The irreducible version (period is 2):



Classification of states

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 - null recurrent if it is recurrent but not positive recurrent.
- transient otherwise, that is, if $\mathbb{P}(\tau_{ii} < \infty) < 1$.

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What about countably infinite chains?

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Given a countable Markov chain with two states i and j in the same strong connected component of the corresponding graph. Then i is transient iff j is.

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Thus with positive probability, a walk starting from i returns to i only finitely many times, a contradiction.

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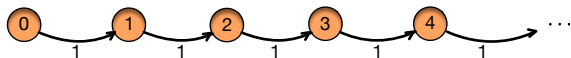
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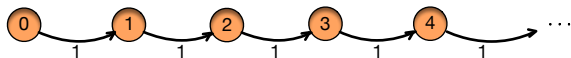
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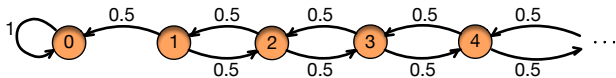
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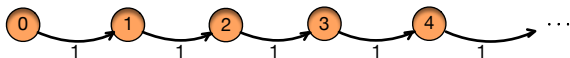
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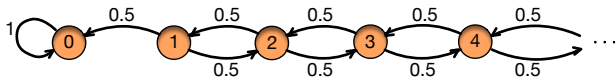
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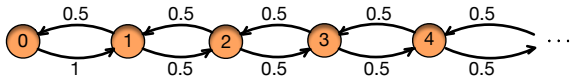
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In the irreducible version, all states are null recurrent (cf. the exercises):



Exercises



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3. The n -lollipop graph consists of a complete graph of $n/2$ vertices with a path of length $n/2$ glued to a vertex. Show that the mean hitting time from any vertex u of the complete graph to the base of the lollipop is at least cubic. Conclude a cubic lower bound for the cover time.



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4. Prove that a connected graph is a two-sided expander iff it is not bipartite. (Hint: recall Hoffman's theorem.)



5. Let $\mathbb{N} \cup \{0\}$ be equipped with the irreducible fair walk structure (third slide): from every positive integer we move to each neighbor with equal probability $1/2$. Show that 0 is a null-recurrent state. (Hint: show that hitting n from 0 has probability $1/n$.)



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6. Classify all states in the previous example, and also in the unfair versions. That is, stepping to the right has probability p . What if $p > 1/2$? And if $p < 1/2$?



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7. Prove the proposition about communicating recurrent states.